

## EE/Ma 127c Error-Correcting Codes - Homework Assignment 2

Ling Li, ling@cs.caltech.edu

April 19, 2001

**2.1** The trellis graph shows

$$W_1 = [1 \ 2], W_2 = \begin{bmatrix} 4 & 0 & 2 \\ 2 & 8 & 0 \end{bmatrix}, W_3 = \begin{bmatrix} 4 & 0 & 2 \\ 8 & 1 & 2 \\ 0 & 1 & 4 \end{bmatrix}, W_4 = \begin{bmatrix} 2 & 0 \\ 4 & 2 \\ 0 & 1 \end{bmatrix}, W_5 = \begin{bmatrix} 1 \\ 2 \end{bmatrix}.$$

(a) From  $\alpha_0 = 1$ ,  $\alpha_i = \alpha_{i-1}W_i$ , for  $i = 1, \dots, 5$ , we get

$$\alpha_0 = 1, \alpha_1 = [1 \ 2], \alpha_2 = [8 \ 16 \ 2], \alpha_3 = [160 \ 18 \ 56], \alpha_4 = [392 \ 92], \alpha_5 = 576;$$

from  $\beta_5 = 1$ ,  $\beta_{i-1} = W_i\beta_i$ , for  $i = 1, \dots, 5$ , we get

$$\beta_0 = 576, \beta_1 = \begin{bmatrix} 80 \\ 248 \end{bmatrix}, \beta_2 = \begin{bmatrix} 12 \\ 28 \\ 16 \end{bmatrix}, \beta_3 = \begin{bmatrix} 2 \\ 8 \\ 2 \end{bmatrix}, \beta_4 = \begin{bmatrix} 1 \\ 2 \end{bmatrix}, \beta_5 = 1.$$

(b) For  $i = 0, 1, \dots, 5$ , and  $v \in V_i$ , let  $\mu_i$  be a column vector of dimension  $q_i$  (same as  $\beta_i$ ), and  $\mu_i(v)$  be the flow from  $A$  to  $B$  through  $v$ . Then  $\mu_i(v) = \mu_v(A, B) = \alpha_i(v)\beta_i(v)$ .

$$\mu_0 = 576, \mu_1 = \begin{bmatrix} 80 \\ 496 \end{bmatrix}, \mu_2 = \begin{bmatrix} 96 \\ 448 \\ 32 \end{bmatrix}, \mu_3 = \begin{bmatrix} 320 \\ 144 \\ 112 \end{bmatrix}, \mu_4 = \begin{bmatrix} 392 \\ 184 \end{bmatrix}, \mu_5 = 576.$$

For  $i = 1, \dots, 5$ , and edge  $e = (u, v)$  where  $u \in V_{i-1}$  and  $v \in V_i$ , we have  $\mu_e(A, B) = \alpha_{i-1}(u)w(e)\beta_i(v)$ . To avoid plotting the trellis graph, let  $\nu_i$  be a  $q_{i-1} \times q_i$  matrix and  $\nu_i(u, v)$  be the flow  $\mu_e(A, B)$ . Thus we have

$$\nu_1 = [80 \ 496], \nu_2 = \begin{bmatrix} 48 & 0 & 32 \\ 48 & 448 & 0 \end{bmatrix}, \nu_3 = \begin{bmatrix} 64 & 0 & 32 \\ 256 & 128 & 64 \\ 0 & 16 & 16 \end{bmatrix}, \nu_4 = \begin{bmatrix} 320 & 0 \\ 72 & 72 \\ 0 & 112 \end{bmatrix}, \nu_5 = \begin{bmatrix} 392 \\ 184 \end{bmatrix}.$$

We can verify that  $\mu_i(v)$  is the same the sum of the  $v^{\text{th}}$  column of  $\nu_i$ .

(c) Take the log (base 2) of  $W_i$ , we get  $W'_i$ :

$$W'_1 = [0 \ 1], W'_2 = \begin{bmatrix} 2 & -\infty & 1 \\ 1 & 3 & -\infty \end{bmatrix}, W'_3 = \begin{bmatrix} 2 & -\infty & 1 \\ 3 & 0 & 1 \\ -\infty & 0 & 2 \end{bmatrix}, W'_4 = \begin{bmatrix} 1 & -\infty \\ 2 & 1 \\ -\infty & 0 \end{bmatrix}, W'_5 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}.$$

Using the “log” forward-backward algorithm with max-sum, we get  $\alpha'_i$  and  $\beta'_i$  below:

$$\alpha'_0 = 0, \alpha'_1 = [0 \ 1], \alpha'_2 = [2 \ 4 \ 1], \alpha'_3 = [7 \ 4 \ 5], \alpha'_4 = [8 \ 5], \alpha'_5 = 8;$$

$$\beta'_0 = 8, \beta'_1 = \begin{bmatrix} 5 \\ 7 \end{bmatrix}, \beta'_2 = \begin{bmatrix} 3 \\ 4 \\ 3 \end{bmatrix}, \beta'_3 = \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix}, \beta'_4 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \beta'_5 = 0.$$

**2.2** Let  $g(\Delta)$  denote the approximation function of  $f(\Delta)$ , which takes the form ( $\Delta_0 = 0, \Delta_4 = \infty$ )

$$g(\Delta) = \begin{cases} y_1, & \Delta_0 \leq \Delta < \Delta_1; \\ y_2, & \Delta_1 \leq \Delta < \Delta_2; \\ y_3, & \Delta_2 \leq \Delta < \Delta_3; \\ y_4, & \Delta_3 \leq \Delta < \Delta_4. \end{cases}$$

The problem is: define an error measure between  $f$  and  $g$  approximation, denoted by  $E(f, g)$ , and find  $g$  that minimizes  $E(f, g)$ .

(a) Define the error measure as the maximum discrepancy between  $f$  and  $g$ :

$$E(f, g) = \sup_{\Delta \geq \Delta_0} |f(\Delta) - g(\Delta)|.$$

(We should use sup instead of max.) Since we know  $f(\Delta)$  is a continuous and monotonously decreasing function of  $\Delta$ , we get for  $i = 1, \dots, 4$ ,

$$\begin{aligned} \sup_{\Delta \in [\Delta_{i-1}, \Delta_i]} |f(\Delta) - g(\Delta)| &= \sup_{\Delta \in [\Delta_{i-1}, \Delta_i]} |f(\Delta) - y_i| \\ &= \max(|f(\Delta_{i-1}) - y_i|, |f(\Delta_i) - y_i|) \\ &\geq \frac{f(\Delta_{i-1}) - f(\Delta_i)}{2}, \end{aligned}$$

with equality iff

$$y_i = \frac{f(\Delta_{i-1}) + f(\Delta_i)}{2}. \quad (1)$$

Thus

$$\begin{aligned} E(f, g) = \sup_{\Delta \geq \Delta_0} |f(\Delta) - g(\Delta)| &= \max_{i=1}^4 \sup_{\Delta \in [\Delta_{i-1}, \Delta_i]} |f(\Delta) - g(\Delta)| \\ &\geq \frac{1}{4} \sum_{i=1}^4 \sup_{\Delta \in [\Delta_{i-1}, \Delta_i]} |f(\Delta) - g(\Delta)| \\ &= \frac{1}{8} f(0) = \frac{\ln 2}{8}, \end{aligned}$$

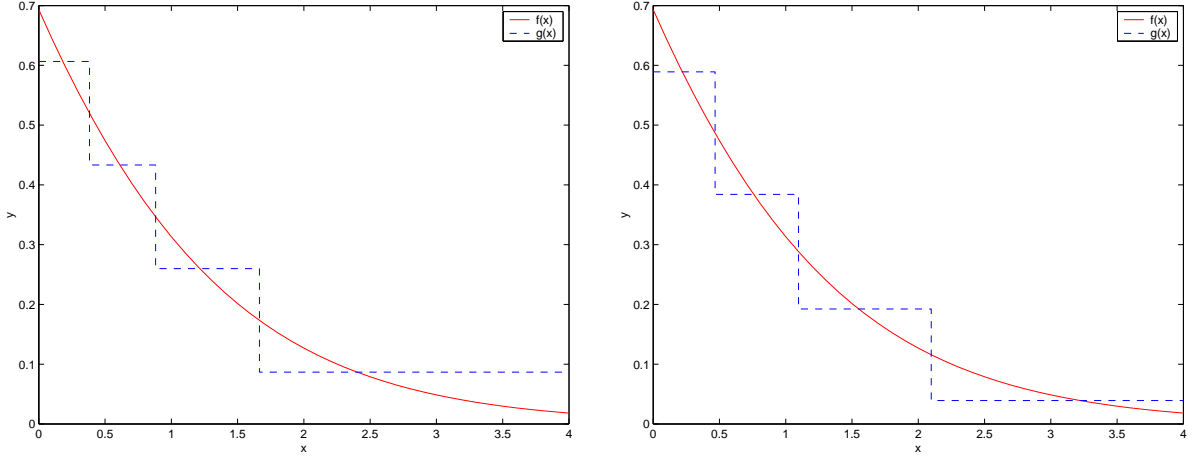
with equality iff (1) holds for all  $i = 1, \dots, 4$ . To minimize  $E(f, g)$ ,  $f(\Delta_i)$  should be  $\frac{4-i}{4} f(0)$ , that is,

$$\Delta_i = f^{-1} \left( \frac{4-i}{4} \ln 2 \right) = -\ln \left( 2^{1-\frac{i}{4}} - 1 \right),$$

and

$$y_i = \frac{9-2i}{8} f(0) = \frac{9-2i}{8} \ln 2.$$

Figure 1(a) shows both  $f$  and  $g$ .



(a) with maximum distance error

(b) with mean square error and  $p = 1/5$ .

Figure 1:  $f(x) = \ln(1 + e^{-x})$  and its approximator  $g(x)$ .

(b) With some predefined probability  $p(x)$ , define the error measure as

$$E(f, g) = \int_0^\infty p(x) (f(x) - g(x))^2 dx = \sum_{i=1}^4 \int_{\Delta_{i-1}}^{\Delta_i} p(x) (f(x) - y_i)^2 dx.$$

To minimize  $E(f, g)$ , we have for  $i = 1, \dots, 4$  and  $j = 1, \dots, 3$ ,

$$\frac{\partial E}{\partial y_i} = 2 \int_{\Delta_{i-1}}^{\Delta_i} p(x) (y_i - f(x)) dx = 0, \quad (2)$$

$$\frac{\partial E}{\partial \Delta_j} = p(\Delta_j) (f(\Delta_j) - y_j)^2 - p(\Delta_j) (f(\Delta_j) - y_{j+1})^2 = 0. \quad (3)$$

From (2), we get

$$y_i = \frac{\int_{\Delta_{i-1}}^{\Delta_i} p(x) f(x) dx}{\int_{\Delta_{i-1}}^{\Delta_i} p(x) dx}.$$

It is reasonable to assume  $p(\Delta_j) \neq 0$ , and  $y_j > f(\Delta_j) > y_{j+1}$  since  $f(\Delta)$  is a decreasing function of  $\Delta$ . Then from (3), we get

$$2f(\Delta_j) = y_j + y_{j+1}.$$

We can not get a closed form for  $y_i$  and  $\Delta_i$ . However, we can use numerical techniques to get a numerical solution. For example, if

$$p(x) = \begin{cases} \frac{1}{5}, & 0 \leq x \leq 5; \\ 0, & \text{otherwise} \end{cases}$$

(since  $f(x) < 0.007$  when  $x > 5$ , we might omit them) I got

$$y_1 \approx 0.589135, y_2 \approx 0.384069, y_3 \approx 0.192328, y_4 \approx 0.039135, \\ \Delta_1 \approx 0.467161, \Delta_2 \approx 1.096547, \Delta_3 \approx 2.098058.$$

Figure 1(b) shows that this  $g$  gives more attention to larger  $x$  than the previous one.

## 2.1 Commutative semiring.

- (a) Below I list the equations of the distributive law in every case. It is really straightforward that they hold.

**sum-product**  $(a \cdot b) + (a \cdot c) = a \cdot (b + c)$ .

**min-product**  $a > 0$ .  $\min(a \cdot b, a \cdot c) = a \cdot \min(b, c)$ .

**max-product**  $a > 0$ .  $\max(a \cdot b, a \cdot c) = a \cdot \max(b, c)$ .

**min-sum**  $\min(a + b, a + c) = a + \min(b, c)$ .

**max-sum**  $\max(a + b, a + c) = a + \max(b, c)$ .

- (b) From mP to MP, we can use the mapping  $x \mapsto x^{-1}$ . That is

$$\frac{\text{mP} \mid \begin{array}{cccccc} x & (0, \infty] & \infty & 1 & a \cdot b & \min(a, b) \end{array}}{\text{MP} \mid \begin{array}{cccccc} x^{-1} & [0, \infty) & 0 & 1 & a^{-1} \cdot b^{-1} & \max(a^{-1}, b^{-1}) \end{array}}$$

From MP to mS, we can use the mapping  $x \mapsto -\log x$ .

$$\frac{\text{MP} \mid \begin{array}{cccccc} x & [0, \infty) & 0 & 1 & a \cdot b & \max(a, b) \end{array}}{\text{mS} \mid \begin{array}{cccccc} -\log x & (-\infty, \infty] & \infty & 0 & (-\log a) + (-\log b) & \min(-\log a, -\log b) \end{array}}$$

From mS to MS, we can use the mapping  $x \mapsto -x$ .

$$\frac{\text{mS} \mid \begin{array}{cccccc} x & (-\infty, \infty] & \infty & 0 & a + b & \min(a, b) \end{array}}{\text{MS} \mid \begin{array}{cccccc} -x & [-\infty, \infty) & -\infty & 0 & (-a) + (-b) & \max(-a, -b) \end{array}}$$

Then by inversing or composing mappings, we get the table below:

$$\begin{array}{c} \text{mP} \quad \text{MP} \quad \text{mS} \quad \text{MS} \\ \left( \begin{array}{cccc} * & x^{-1} & \log x & -\log x \\ x^{-1} & * & -\log x & \log x \\ e^x & e^{-x} & * & -x \\ e^{-x} & e^x & -x & * \end{array} \right) \end{array}$$